

Boolean Algebra

Def A boolean algebra $(\beta, +, \cdot, -, 0, 1)$ is a set, β , with two binary operators, $+$, \cdot , one unary operator, $-$, and two elements $0, 1 \in \beta$ with the following properties:

1. There exists at least two elements $a, b \in \beta$ with $a \neq b$.
2. $\forall a, b \in \beta, a + b \in \beta, a \cdot b \in \beta$ (Closure)
3. $\forall a, b \in \beta, a + b = b + a, a \cdot b = b \cdot a$ (Commutative)
4. $\forall a \in \beta, a + 0 = a, a \cdot 1 = a$ (Identity)
5. $\forall a, b, c \in \beta$
 $a \cdot (b + c) = a \cdot b + a \cdot c$ (Distributive)
 $a + (b \cdot c) = (a + b) \cdot (a + c)$
6. $\forall a \in \beta, \exists \bar{a} \in \beta$ with $a + \bar{a} = 1$ (Complement)
 $a \cdot \bar{a} = 0$
7. $\forall a, b, c \in \beta$
 $a + (b + c) = (a + b) + c$ (Associative)
 $a \cdot (b \cdot c) = (a \cdot b) \cdot c$

Boolean algebras we have studied:

β	$\{T, F\}$	$P(X)$
$+$	\vee	\cup
\cdot	\wedge	\cap
$-$	\neg	$\overline{(\quad)}$
0	F	\emptyset
1	T	\mathcal{U}

Principle of Duality

Every property or theorem has a dual property or theorem (interchange $+$ and \cdot , interchange 0 and 1).

Theorems for all boolean algebras

Thm. $\forall a \in \beta \quad a+a=a$ Dual $(a \cdot a=a)$

$$\begin{aligned} a+a &= (a+a) \cdot 1 && \text{Identity} \\ &= (a+a) \cdot (a+\bar{a}) && \text{Complement} \\ &= a + a \cdot \bar{a} && \text{Distributive} \\ &= a + 0 && \text{Complement} \\ &= a && \text{Identity} \end{aligned}$$

Thm. $\forall a \in \beta \quad a+1=a$ Dual $(a \cdot 0=0)$

$$\begin{aligned} a+1 &= (a+1) \cdot 1 && \text{Identity} \\ &= (a+1) \cdot (a+\bar{a}) && \text{Complement} \\ &= a + 1 \cdot \bar{a} && \text{Distributive} \\ &= a + \bar{a} && \text{Identity} \\ &= 1 && \text{Complement} \end{aligned}$$

Thm. $\forall a \in \beta \quad a+a \cdot b=a$ Dual $a \cdot (a+b)=a$

$$\begin{aligned} a+a \cdot b &= a \cdot 1 + a \cdot b && \text{Identity} \\ &= a \cdot (1+b) && \text{Distributive} \\ &= a \cdot 1 && \text{Thm } a+1=1 \\ &= a && \text{Identity} \end{aligned}$$

Thm. $\forall a, b \in \beta \quad a + \bar{a} \cdot b = a+b$ Dual $a \cdot (\bar{a}+b) = a \cdot b$

$$\begin{aligned} a + \bar{a} \cdot b &= (a + \bar{a}) \cdot (a+b) && \text{Distributive} \\ &= 1 \cdot (a+b) && \text{Complement} \\ &= a+b && \text{Identity} \end{aligned}$$

Boolean Functions

Consider the boolean algebra
($\{T, F\}, \vee, \wedge, \neg, \top, \perp$) but use the notation
($\{1, 0\}, +, \cdot, -, 1, 0$) ($\beta = \{1, 0\}$)

Def A boolean function of n variables is a function $f: \beta^n \rightarrow \beta$ where $\beta = \{1, 0\}$

$$\beta^n = \{(x_1, x_2, \dots, x_n) \mid x_i \in \beta, 1 \leq i \leq n\}$$

Ex $f: \beta^3 \rightarrow \beta$ $f(x, y, z) = x \cdot y + \bar{z}$

x	y	z	f
1	1	1	1
1	1	0	1
1	0	1	0
1	0	0	1
0	1	1	0
0	1	0	1
0	0	1	0
0	0	0	1

Different boolean expressions determine the same boolean functions:

$$f(x_1, x_2, x_3) = x_1 \cdot (x_2 + x_3) = x_1 \cdot x_2 + x_1 \cdot x_3$$

$$f(x_1, x_2) = \overline{x_1 \cdot x_2} = \bar{x}_1 + \bar{x}_2$$

Standardized Form for boolean functions

Def A literal is a boolean variable or its

complement: $y = x$ or $y = \bar{x}$

A minterm of the variables x_1, x_2, \dots, x_n is a boolean product of literals: $y_1 \cdot y_2 \cdot \dots \cdot y_n$

EX Two variables x_1, x_2

$y_1 \cdot y_2$ y_1 literal for x_1 y_2 literal for x_2

$x_1 \cdot x_2$

$x_1 \cdot \bar{x}_2$

$\bar{x}_1 \cdot \bar{x}_2$

$\bar{x}_1 \cdot x_2$

4 minterms for 2 variables

2^n minterms for n variables

There is a bijection between n tuples and minterms of n variables

$$(j_1, j_2, \dots, j_n) \rightarrow y_1 \cdot y_2 \cdot \dots \cdot y_n$$

If $j_i = 1$ then $y_i = x_i$

$j_i = 0$ then $y_i = \bar{x}_i$

EX. $(1, 0, 0, 1) \rightarrow x_1 \cdot \bar{x}_2 \cdot \bar{x}_3 \cdot x_4$

$$(0, 0, 1, 1, 0) \rightarrow \bar{x}_1 \cdot \bar{x}_2 \cdot x_3 \cdot x_4 \cdot \bar{x}_5$$

Note: $\bar{x}_1 \cdot \bar{x}_2 \cdot x_3 \cdot x_4 \cdot \bar{x}_5 = 1$ iff $x_1 = 0, x_2 = 0, x_3 = x_4 = 1, x_5 = 0$

Def. Standardized form for boolean functions is called Sum of Products Form or Disjunctive Normal Form (DNF) and is the representation of the function as a sum of minterms.
(Unique up to the ordering of minterms)

EX. Find the sum of products form for $f(x, y, z) = (x+y) \cdot z$

① Calculate the truth table

x	y	z	$(x+y) \cdot \bar{z}$
1	1	1	0
1	1	0	1
1	0	1	0
1	0	0	1
0	1	1	0
0	1	0	1
0	0	1	0
0	0	0	0

② Identify the n-tuples at which the function is 1

③ Form the minterms corresponding to these n-tuples and sum.

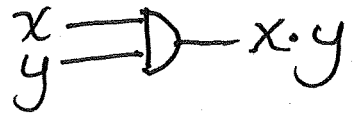
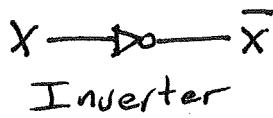
$$f(x, y, z) = x \cdot y \cdot \bar{z} + x \cdot \bar{y} \cdot \bar{z} + \bar{x} \cdot y \cdot \bar{z}$$

Sum of products form (DNF)

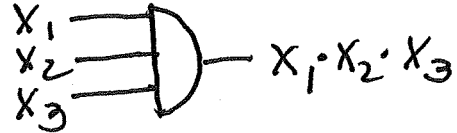
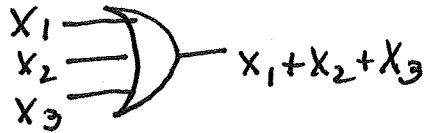
Note: we have also proved that every boolean function can be accomplished using disjunction, conjunction and negation.

Circuits - physical application of boolean functions

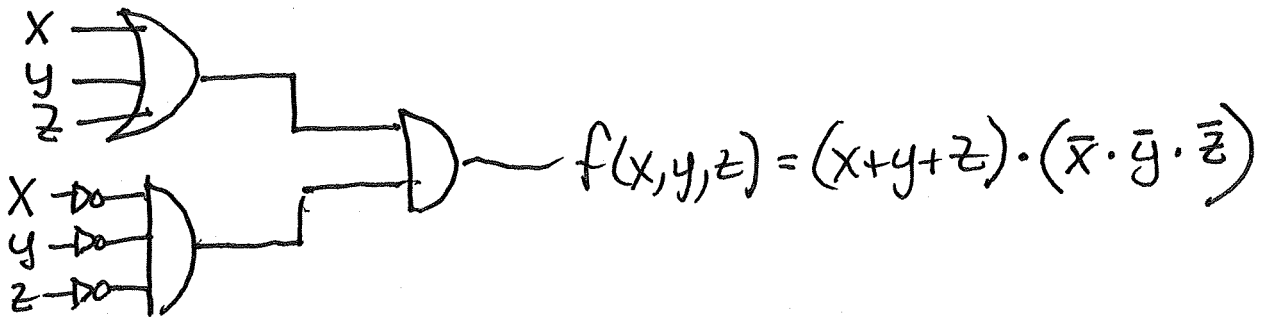
Circuits made up of gates:



Allow in design:



Represent a circuit with a boolean function:



Minimization of Circuits

Represent a circuit by a boolean function.
Get an equivalent boolean function with a minimal number of terms in a sum of product expression.

EX. (Algebraic method)

$$\begin{aligned}
 f(x_1, x_2, x_3) &= \overline{x_1} \cdot x_2 \cdot \overline{x_3} + \overline{x_1} \cdot \overline{x_2} \cdot \overline{x_3} + x_1 \cdot \overline{x_2} \cdot \overline{x_3} + x_1 \cdot x_2 \cdot x_3 + x_1 \cdot \overline{x_2} \cdot x_3 \\
 &= \overline{x_1} \cdot \overline{x_3} (x_2 + \overline{x_2}) + x_1 \cdot \overline{x_2} (\overline{x_3} + x_3) + x_1 x_3 (x_2 + \overline{x_2}) \\
 &= \overline{x_1} \cdot \overline{x_3} (1) + x_1 \cdot \overline{x_2} (1) + x_1 x_3 (1) \\
 &= \overline{x_1} \cdot \overline{x_3} + x_1 \overline{x_2} + x_1 x_3
 \end{aligned}$$

Karnaugh maps - visual aid to minimize circuits.

	y	\bar{y}
x		
\bar{x}		

Every cell represents a minterm.
 If the minterm appears in the sum of products form, put a 1 in the box.
 Adjacent cells represent minterms which differ in only one variable, so adjacent cells can be collapsed.

- Combine 2 adjacent cells - eliminate 1 variable
- 4 adjacent cells - eliminate 2 variables
- 8 adjacent cells - eliminate 3 variables

Examples (Taken circuits, represented them by boolean functions, put in sum of products form, then filled in the Karnaugh map)

	y	\bar{y}	
x	1		
\bar{x}	1		

$f(x,y) = y$

	y	\bar{y}	
x		1	
\bar{x}	1	1	

$f(x,y) = \bar{x} + \bar{y}$

Three variables:

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
x	*			*
\bar{x}	**			**

(Right adjacent to left, *, **)

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
x	1			
\bar{x}	1			

$f(x,y,z) = y \cdot z$

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
x	1	1		
\bar{x}	1	1		

$f(x,y,z) = y$ (rather than $yz + y\bar{z}$)

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
x	1	1	1	1
\bar{x}				1

$$f(x,y,z) = x + \bar{y} \cdot z$$

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
x	1	1		1
\bar{x}	1	1		

$$f(x,y,z) = y + z$$

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
x	1	1		
\bar{x}			1	1

$$f(x,y,z) = xy + \bar{x} \cdot \bar{y}$$

Four variables

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
wx	*			*
$w\bar{x}$				
$\bar{w}x$				
$\bar{w}\bar{x}$	*			*

* adjacent cells

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
wx	1	1		1
$w\bar{x}$	1	1		
$\bar{w}x$			1	1
$\bar{w}\bar{x}$				

$$f(x,y,z,w) = wy + wxz + \bar{w}\bar{x}\bar{y}$$

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
wx	1			1
$w\bar{x}$		1	1	
$\bar{w}x$	1			1
$\bar{w}\bar{x}$				

$$f(x,y,z,w) = xz + \bar{x} \cdot \bar{z}$$

	yz	$y\bar{z}$	$\bar{y}z$	$\bar{y}\bar{z}$
wx	1			1
$w\bar{x}$	1			1
$\bar{w}x$	1	1	1	1
$\bar{w}\bar{x}$	1	1	1	1

$$f(x,y,z,w) = \bar{w} + z$$